

Twice-Ramanujan Sparsifiers

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Aim: . Compute k-sparsifier H of graph G .

$$\Rightarrow x^T L_G x \leq x^T \underbrace{L_H}_{\Rightarrow \text{Laplacian}} x \leq k \cdot x^T L_G x$$

s.t. number of edges is less in H .

W
 $D = \text{diag}$
 $D_{ii} = \sum_{a:i} w_i$

* Thm: $\# G = (V, E, w)$: Undirected weighted graph
 $\# H = (V, E, \tilde{w})$: $\lceil d(n-1) \rceil$ edges

$$x^T L_G x \leq x^T L_H x \leq \underbrace{\left(\frac{d+1+2\sqrt{d}}{d+1-2\sqrt{d}} \right)}_k \cdot x^T L_G x$$

complete $\Rightarrow d(n-1)$
 Ramanujan sparsifiers.

$$1 + \frac{4\sqrt{d}}{d+1-2\sqrt{d}}$$

$$1 + \frac{4}{d-2}$$

Preliminaries

- $L_G = D - W$

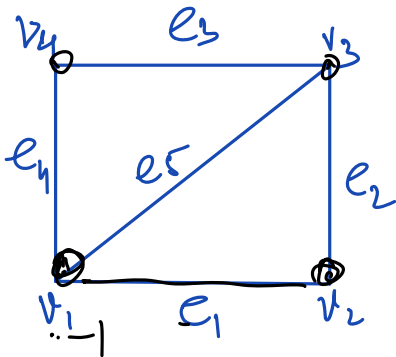


- $x^T L_G x = \sum_{(u,v) \in E} w_{u,v} (x_u - x_v)^2$

$$x = \begin{pmatrix} -1 \\ -1 \\ \vdots \\ -1 \end{pmatrix}$$

- Incidence Matrix : $B_{m \times n}$ (m -edges, n -vertices)

$$B(e, v) = \begin{cases} 1 & \text{if } v \text{ is } e\text{'s head} \\ -1 & \text{if } v \text{ is } e\text{'s tail} \\ 0 & \text{otherwise} \end{cases}$$



$$B = \begin{matrix} & \begin{matrix} v_1 & v_2 & v_3 & v_4 \end{matrix} \\ \begin{matrix} \cdot e_1 \\ \cdot e_2 \\ \cdot e_3 \\ \cdot e_4 \\ \cdot e_5 \end{matrix} & \begin{bmatrix} -1 & 1 & 0 & 0 \\ 0 & -1 & 1 & 0 \\ 0 & 0 & -1 & 1 \\ 1 & 0 & 0 & -1 \\ -1 & 0 & 1 & 0 \end{bmatrix} \end{matrix}$$

$$W_{m \times m} = \text{diag}(w_e)$$

Note: $x^T L_G x = x^T B^T W B x = \|W^{\frac{1}{2}} B x\|_2^2$

$$= \sum_{(u,v) \in E} w_{uv} (x_u - x_v)^2, \quad \forall x \in \mathbb{R}^n$$

- G connected $\Leftrightarrow \underline{\text{Ker}(L)} = \underline{\text{Ker}(W^{\frac{1}{2}} B)} = \text{span}(\underline{\mathbb{1}})$

Rank-1 Updates

* (Sherman-Morrison): If A is non singular $n \times n$,
v is a vector, then
 \Rightarrow

$$\underbrace{(A + vv^T)^{-1}}_{\leftarrow \quad \rightarrow} = A^{-1} - \frac{A^{-1} v v^T A^{-1}}{1 + v^T A^{-1} v} :$$

\Downarrow

* (Matrix Determinant Lemma):

$$\underbrace{\det(A + vv^T)}_{\leftarrow \quad \rightarrow} = \underbrace{\det(A)}_{\leftarrow \quad \rightarrow} \underbrace{(1 + v^T A^{-1} v)}_{\leftarrow \quad \rightarrow}$$
$$A (I + A^{-1} v v^T)$$

Proof Sketch: Factor $A + v v^T = A (\underbrace{I + A^{-1} v v^T}_{\text{rank 1 update}})$

Now consider $(\underline{I + u v^T})$:

$$\underline{(I + u v^T)u} = \underline{(1 + v^T u)u}$$

\Downarrow
eigenvectors = u ,
eigenvals = $(1 + u^T v)$, $\text{Span}(u^\perp)$
 1

$(I + v v^T)^{-1} \rightarrow$ Same eigenvectors: v , $\text{Span}(v^\perp)$
eigenvals: $\frac{1}{(1 + \|v\|^2)}$, 1

\therefore Any matrix of form $(\underline{I + g v v^T})$ must have all 1's eigenval except one $1 + g \|v\|^2$.

$$\therefore \frac{(I + vv^T)(I + vv^T)^{-1}}{\downarrow} = I$$

$$\Rightarrow (I + vv^T)(I + gvv^T) = I$$

$$\Rightarrow I + vv^T + gvv^T + v g \|v\|^2 v^T = I$$

$$\therefore I + v(1 + g + g\|v\|^2)v^T = I$$

$$\Rightarrow (1 + g + g\|v\|^2) = 0$$

$$\therefore g = \frac{-1}{1 + \|v\|^2}$$

$$\therefore \frac{(I + vv^T)^2}{\underbrace{\quad}_{g}} = I + v \left(\frac{-1}{1 + v^T v} \right) v^T$$

□

Main Result

* Thm: Suppose $d > 1$, $v_1, \dots, v_m \in \mathbb{R}^n$ s.t.

$$\sum_i v_i v_i^T = I_n \quad m > n$$

Then \exists scalars $s_i \geq 0$ with $|\{i: s_i \neq 0\}| \leq dn$
with $\sum_{i \in \{m\}} s_i v_i v_i^T = I_n$

$$I_n \preceq \sum_i s_i v_i v_i^T \preceq \left(\frac{d+1+2\sqrt{d}}{d+1-2\sqrt{d}} \right) I_n$$

↔

Theorem \Rightarrow Sparsification

* Proof of Thm \Rightarrow Sparsification:

wlog G is connected. $L_G = B^T W B$. Fix $d > 1$.
Restrict attention to

$$\underline{\text{Im}(L_G) \equiv \mathbb{R}^{n-1}}.$$

Apply Thm to columns $\{v_i\}_{i \in m}$ of

$$\underline{V_{n \times m}} = \underline{(L_G^+)^{\frac{1}{2}}} B^T W^{\frac{1}{2}}.$$

Note: $\underline{\sum_i v_i v_i^T} = \underline{V V^T} = (L_G^+)^{\frac{1}{2}} \underbrace{B^T W B}_{L_G} (L_G^+)^{\frac{1}{2}}$
 $= \underline{(L_G^+)^{\frac{1}{2}} (L_G) (L_G^+)^{\frac{1}{2}}} = \underline{\underline{I_{\text{Im}(L_G)}}}$

* Now, we construct sparser graph from theorem.

let $S = \text{diag}(s_i)$.

$$L_H = \underline{B^T} \underline{W^{\frac{1}{2}}} \underline{S} \underline{W^{\frac{1}{2}}} \underline{B} \Rightarrow \text{sub graph of } G \text{ with} \\ \text{weights } (\tilde{w}_i = w_i s_i)$$

Also:

$$\underline{I_{\text{Im}(L_G)}} \preceq \sum s_i v_i v_i^T = \underline{V S V^T} \preceq \kappa \cdot \underline{I_{\text{Im}(L_G)}}$$

\therefore we get

$$1 \leq \frac{y^T \underline{V S V^T} y}{y^T y} \leq \kappa$$

$$\underline{x^T L_G x \leq x^T L_H x \leq \kappa \cdot x^T L_G x}$$

$$\Rightarrow 1 \leq \frac{y^T (L_G^+)^{\frac{1}{2}} L_H (L_G^+)^{\frac{1}{2}} y}{y^T y} \leq \kappa \quad (\forall y \in \text{Im}(L_G))$$

$$(\because V = (L_G^+)^{\frac{1}{2}} B^T W^{\frac{1}{2}})$$

$$\Rightarrow 1 \leq \frac{\underline{x^T L_G^{\frac{1}{2}} (L_G^+)^{\frac{1}{2}} L_H (L_G^+)^{\frac{1}{2}} L_G^{\frac{1}{2}} x}}{\underline{x^T L_G^{\frac{1}{2}} L_G^{\frac{1}{2}} x}}$$

$$(\forall x \perp \mathbb{1})$$

$$Lx = y$$

$$\Rightarrow \boxed{1 \leq \frac{x^T L_n x}{x^T L_G x} \leq \rho} \quad (\forall x \perp \mathbb{1})$$

\therefore Theorem \Rightarrow Sparsification



(Recall)

* Thm: Suppose $d > 1$, $v_1, \dots, v_m \in \mathbb{R}^n$ s.t.

$$\sum_i v_i v_i^T = I_{n \times n}.$$


Then \exists scalars $s_i \geq 0$ with $|\{i: s_i \neq 0\}| \leq dn$
with

$$I_n \preceq \sum_i s_i v_i v_i^T \preceq \left(\frac{d+1+2\sqrt{d}}{d+1-2\sqrt{d}} \right) I_n$$

Proof Intuition

- Eigenvalues of A interlace eigenvalues of $A+vv^T$.
- We can determine the new eigenvalues:

$$p_{A+vv^T}(\lambda) = \underbrace{p_A(\lambda)}_{\det(\lambda I - A - vv^T)} \left(1 - \sum_j \frac{\langle v, u_j \rangle^2}{\lambda - \lambda_j} \right)$$




 $\det(A+vv^T) = \det(\lambda I) \cdot (I^{-1}v)^T$

The $p_{A+vv^T}(\lambda)$ has 2 kinds of zeroes λ :

- Those for which $p_A(\lambda) = 0$, i.e. $v \perp u_j$
- Those for which $p_A(\lambda) \neq 0$ & $p_{A+vv^T}(\lambda) = 0$

$$\Rightarrow f(\lambda) := \left(1 - \frac{1}{m} \sum_j \frac{\langle v, u_j \rangle^2}{\lambda - \lambda_j} \right) = 0$$



⇒ These eig have "moved", & interlace old eigenvals.

• Now suppose we add random vec.: $A + v_i v_i^T$

$$\mathbb{E}_v \langle v, u_j \rangle^2 = \frac{1}{m} \sum_i \langle v_i, u_j \rangle^2 = \frac{1}{m} u_j^T \left(\sum_i v_i v_i^T \right) u_j$$

$$\frac{\sum v_i v_i^T}{m} = I$$

$$= \frac{\|u_j\|^2}{m} = \frac{1}{m}$$

∴ Adding this "average" vec. v_i gives:

$$P_{A+v_i v_i^T}(\lambda) = P_A(\lambda) \left(1 - \frac{1}{m} \frac{1}{\lambda - \lambda_i} \right) = P_A(\lambda) - \frac{1}{m} P_A'(\lambda)$$

$$\therefore P_A' = \sum_j \prod_{i \neq j} (\lambda - \lambda_i) \quad P_A(\lambda) = \prod_i (\lambda - \lambda_i)$$

Start with $A=0 \Rightarrow P_A(\lambda) = \lambda^m$

$$P_{A+vv^T}(\lambda) = \left(P_A(\lambda) - \frac{1}{m} P_A'(\lambda) \right) - \frac{1}{m} P_{A+vv^T}(\lambda)$$

∴ ⇒ Laguerre Polynomial

"Well-known" for Laguerre that after d iterations:

$$\frac{\text{Largest zero}}{\text{Smallest zero}} = \frac{d+1+2\sqrt{d}}{d+1-2\sqrt{d}}$$

(Recall)

* Inv: Suppose $d > 1$, $v_1, \dots, v_m \in \mathbb{R}^n$ s.t.

$$\sum_i v_i v_i^T = I_{n \times n}.$$

Then \exists scalars $s_i \geq 0$ with $|\{i: s_i \neq 0\}| \leq d$ with

$$I_n \preceq \sum_i s_i v_i v_i^T \preceq \frac{d+1+2\sqrt{d}}{d+1-2\sqrt{d}} I_n$$

Proof by Barrier Functions

* Def: $\underline{u}, \underline{l} \in \mathbb{R}$, A symmetric $\lambda_1, \dots, \lambda_n$

$$\underline{\Phi}^u(A) := \text{tr}(\underline{uI} - A)^{-1} = \sum_i \frac{1}{\underline{u} - \lambda_i} \quad (\text{Upper Potential})$$

$$\underline{\Phi}^l(A) := \text{tr}(A - \underline{lI})^{-1} = \sum_i \frac{1}{\lambda_i - \underline{l}} \quad (\text{Lower Potential})$$



* Rmk: As long as $A \preceq uI$, $A \succeq lI$, potential measures how "far" eig(A) are from barriers u & l .

* Intuition: Suppose $\Phi^\mu(A) \leq 1$, \checkmark $A \leq \mu I$. Then,

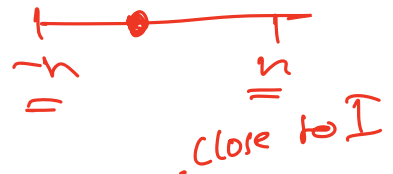
$$\Phi^\mu(A) = \sum_i \frac{1}{\mu - \lambda_i} \leq 1$$



\Rightarrow At most 1 eig at dist 1 from μ .
At most 2 eig at dist 2 from μ .
:

$\Phi^\mu(A)$ measures total "repulsion" of eigenvalues of A from potential μ .

* Steps to prove the theorem:



Iteratively build $\sum_i s_i v_i v_i^T$.

Start from $A^{(0)} = 0$
 $A^{(1)} = A^{(0)} + \tau v_1 v_1^T$

Consider the consts: $u_0, l_0, \underline{\delta}_u, \underline{\delta}_L, \underline{\epsilon}_u, \underline{\epsilon}_L$

(a) Initial Barriers: $u = u_0, l = l_0$, & potentials:

$$\Phi^{u_0}(A^{(0)}) = \epsilon_U$$

$$\Phi^{l_0}(A^{(1)}) = \epsilon_L$$

$$\sum \frac{1}{u_0 - \lambda_i} = \epsilon_U$$

\Rightarrow

$$\sum \frac{1}{u - \lambda_i} = \sum \frac{1}{u + \delta_u - \lambda_i}$$

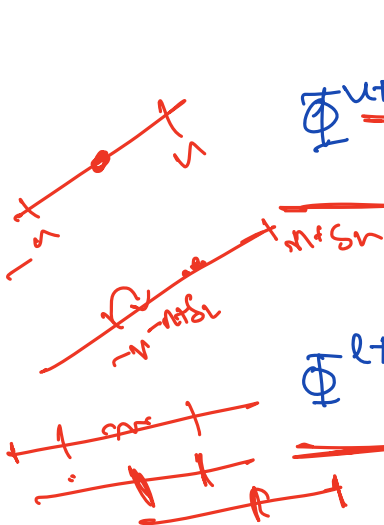
$A^T V U^T$

(b) Each matrix obtained by:

$$A^{(q+1)} = A^{(q)} + t v v^T, \quad v \in \{v_1, \dots, v_m\}$$

$t \geq 0$.

(c) Increment Barriers: $u \rightarrow u + \delta u$ $l \rightarrow l + \delta l$:



$$\Phi^{u+\delta u}(A^{(q+1)}) \leq \Phi^u(A^{(q)}) \leq \delta u,$$

where $u = u_0 + q \delta u$

$$\Phi^{l+\delta l}(A^{(q+1)}) \leq \Phi^l(A^{(q)}) \leq \delta l$$

where $l = l_0 + q \delta l$

(d) No eig jumps across barriers:

$$\lambda_{\max}(A^{(q)}) \leq u_0 + q \delta u \quad \& \quad \lambda_{\min}(A^{(q)}) \geq l_0 + q \delta l.$$

Combining (a), (b), (c), (d), we get: for $Q = dn$

$$\frac{\lambda_{\max}(A^{(Q)})}{\lambda_{\min}(A^{(Q)})} \leq \frac{20 + dn \delta_u}{20 + dn \delta_L} = \frac{d+1 + 2\sqrt{d}}{d+1 - 2\sqrt{d}}$$

by construction: $A^{(Q)}$ is weighted sum of
at most dn vectors.

$$\pm v_i v_i^T$$

* Lemma 1: (Upper Barrier Shift)

Shifting u to $u + \delta u$, we observe

$$\underline{\Phi}^u(A) > \underline{\Phi}^{u+\delta u}(A).$$

\therefore Can add $(A + tVV^T)$ s.t.

$$\underline{\Phi}^u(A) > \underline{\Phi}^{u+\delta u}(A + tVV^T).$$

$t = ?$
 \Rightarrow

Suppose $\lambda_{\max}(A) \leq u$, v is any vector. If

$$\frac{1}{t} \geq v \left(\frac{((u+\delta u)I - A)^{-2}}{\Phi^u(A) - \Phi^{u+\delta u}(A)} + ((u+\delta u)I - A)^{-1} \right) v^T =: \underline{u}_A(t)$$

$$\underline{\Phi}^{u+\delta u}(A + tVV^T) \leq \underline{\Phi}^u(A) \quad \vee \quad \lambda_{\max}(A + tVV^T) \leq \underline{u + \delta u}.$$

* Proof: let $u' = u + \epsilon v$.

$$\Phi^{u+\epsilon v}(A+t vv^T) = \text{tr} (u'I - A - t vv^T)^{-1}$$

$$= \text{tr} \left((u'I - A)^{-1} + \frac{t (u'I - A)^{-1} vv^T (u'I - A)^{-1}}{1 - t v^T (u'I - A)^{-1} v} \right)$$

$$= \text{tr} (u'I - A)^{-1} + t \cdot \frac{v^T (u'I - A)^{-2} v}{1 - t v^T (u'I - A)^{-1} v}$$

$$= \Phi^u(A) - \left(\Phi^v(A) + \Phi^{u'}(A) \right) + \frac{v^T (u'I - A)^{-2} v}{\underbrace{\frac{1}{t} - v^T (u'I - A)^{-1} v}}$$

∴

$$U_A(v) > v^T (u'I - A)v \quad \& \quad \frac{1}{t} \geq U_A(v)$$

∴ last term is finite

\Rightarrow By choice of $\frac{1}{t}$

$$\frac{v^T (U^T I - A)^{-2} v}{\frac{1}{t} - v^T (U^T I - A)^{-1} v}$$

Last term

$$\frac{1}{t} \geq v^T \left(\frac{((U + \delta U) I - A)^{-2}}{\Phi^U(A) - \Phi^{U + \delta U}(A)} + ((U + \delta U) I - A)^{-1} \right) v^T, \text{ we get}$$

$$\Phi^{U + \delta U}(A) \leq \Phi^U(A).$$

Moreover since last term is finite,

$$\lambda_{\max}(A + t v v^T) < U + \delta U.$$

because if not, then for some true $t' < t$
 $\lambda_{\max}(A + t' v v^T) = U + \delta U$. But at such t'

$\Phi^{U + \delta U}(A + t' v v^T)$ would be as \Rightarrow contradiction



* Lemma 2: (lower Barrier)

Suppose $\lambda_{\min}(A) > l$, $\Phi_l(A) \leq \frac{1}{\delta_L}$, v any vec.
If,

$$0 < \frac{1}{t} \leq v \left(\frac{(A - (l + \delta_L)I)^{-2}}{\Phi^{l+\delta_L}(A) - \Phi^l(A)} - (A - (l + \delta_L)I)^{-1} \right) v =: L_A(v)$$

then,

$$\Phi_{l+\delta_L}(A + t v v^T) \leq \Phi_l(A) \quad \text{and}$$

$$\lambda_{\min}(A + t v v^T) > l + \delta_L$$

$$L_A(v) \leq \frac{1}{t} \leq L_A(v)$$

same proof.

* Lemma 3: (Both Barriers Simultaneously)

If $\lambda_{\max}(A) < u$, $\lambda_{\min}(A) > l$, $\Phi^u(A) \leq \epsilon_u$, $\Phi^l(A) \leq \epsilon_l$,

and $\epsilon_u, \epsilon_l, \delta_u$ & δ_l satisfy:

$$0 \leq \underbrace{\frac{1}{\delta_u} + \epsilon_u} \leq \underbrace{\frac{1}{\delta_l} - \epsilon_l}. \text{ Then } \exists i, t > 0 \text{ s.t.}$$

$$\lambda_A(v_i) \geq \frac{1}{t} \geq \lambda_A(v_i), \quad \lambda_{\max}(A + t v_i v_i^T) < u + \delta_u,$$

$$\lambda_{\min}(A + t v_i v_i^T) > l + \delta_l.$$

* Proof : We will show :

$$\sum_i L_A(v_i) \geq \sum_i U_A(v_i)$$

\Downarrow

$$\exists \text{ vector } v_i \text{ s.t. } L_A(v_i) \geq U_A(v_i)$$

$$\sum v_A(v_i) = \sum_i v_i \left(\frac{((\mu + \delta u)I - A)^{-2}}{\Phi^u(A) - \Phi^{u+\delta u}(A)} + ((\mu + \delta u)I - A)^{-1} \right) v_i^T$$

$$= \frac{((\mu + \delta u)I - A)^{-2}}{\Phi^u(A) - \Phi^{u+\delta u}(A)} \cdot \underbrace{\left(\sum_i v_i v_i^T \right)}_{=I} + ((\mu + \delta u)I - A)^{-1} \cdot \underbrace{\left(\sum_i v_i v_i^T \right)}_{=I}$$

$$= \text{tr} \left(\frac{((\mu + \delta u)I - A)^{-2}}{\Phi^u(A) - \Phi^{u+\delta u}(A)} \right) + \text{tr} \left(((\mu + \delta u)I - A)^{-1} \right)$$

$$= \frac{\sum_i (\mu + \delta u - \lambda_i)^{-2}}{\sum_i (\mu - \lambda_i)^{-1} - \sum_i (\mu + \delta u - \lambda_i)^{-1}} + \Phi^{u+\delta u}(A)$$

$$= \frac{\sum_i (\mu + \delta u - \lambda_i)^{-2}}{\delta u \left(\sum_i (\mu - \lambda_i)^{-1} \cdot (\mu + \delta u - \lambda_i)^{-1} \right)} + \Phi^{u+\delta u}(A)$$

$$\therefore \sum (u - \lambda_i)^{-1} (u + \delta u - \lambda_i)^{-1} \geq \sum_i (u + \delta u - \lambda_i)^{-2}$$

$$\Rightarrow \sum_i U_A(v_i) \leq \frac{1}{\delta_L} + \Phi^{u + \delta u}(A) \leq \frac{1}{\delta_u} + \epsilon_u$$

Similarly,

$$\sum_i L_A(v_i) \geq \frac{1}{\delta_L} - \sum_i (\lambda_i - l)^{-1} = \frac{1}{\delta_L} - \epsilon_L$$

Putting these together

$$\sum_i U_A(v_i) \leq \frac{1}{\delta_u} + \epsilon_u \leq \frac{1}{\delta_L} - \epsilon_L \leq \sum_i L_A(v_i) \quad \square$$

Proof of Theorem

We need to set $\epsilon_u, \epsilon_l, \delta_u, \delta_l$ s.t. Lemma 3 is satisfied.

Recall: Constructed $A^{(q+1)} = A^{(q)} + t v_i v_i^T, t \geq 0$

$$\rightarrow L_{A^{(q)}}(v_i) \geq \frac{1}{t} \geq U_{A^{(q)}}(v_i)$$

Take $\delta_l = 1$, $\epsilon_l = \frac{1}{\sqrt{d}}$, $\delta_0 = -n/\epsilon_l$

$\delta_u = \frac{\sqrt{d}+1}{\sqrt{d}-1}$, $\epsilon_u = \frac{\sqrt{d}-1}{d+\sqrt{d}}$, $\delta_0 = n/\epsilon_u$

check that

$$\frac{1}{\delta_u} + \epsilon_u = \frac{1}{\delta_l} - \epsilon_l$$

$$\begin{array}{cc} -n & +n \\ \epsilon_l & \epsilon_u \\ \delta_u + \delta_l & \delta_u + \delta_l \end{array}$$

Initial potentials are: $\Phi_{\frac{n}{\epsilon_0}}(0) = \epsilon_0$
 $\Phi_{\frac{n}{\epsilon_L}}(0) = \epsilon_L$

$$\frac{\lambda_{\max}(A^{(dn)})}{\lambda_{\min}(A^{(dn)})} \leq \frac{\frac{n}{\epsilon_0} + dn \epsilon_0}{-\frac{n}{\epsilon_L} + dn \epsilon_L}$$

$$\approx \frac{\frac{d + \sqrt{d}}{\sqrt{d} - 1} + \frac{d\sqrt{d} + 1}{\sqrt{d} - 1}}{d - \sqrt{d}}$$

$$= \frac{d + 2\sqrt{d} + 1}{d - 2\sqrt{d} + 1}$$

□

The Algorithm:

- Compute vectors $v_i \Rightarrow \mathcal{O}(n^2m)$
- Compute $\left((u+su)I - A\right)^{-1}$, $\left((u+su)I - A\right)^2$ in each iter
 $= \mathcal{O}(n^3)$
- Compute $U_A(v_i) \approx L_A(v_i) = \mathcal{O}(n^2m)$
- Run for (dn) iterations \Rightarrow
 $\mathcal{O}(dn^3m)$.